Overview of the Global Arrays Parallel Software Development Toolkit

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Overview

- Background
- Programming Model
- Core Capabilities
- New Functionality
- Applications
- Summary
Distributed Data:

Data is explicitly associated with each processor, accessing data requires specifying the location of the data on the processor and the processor itself.

Data locality is explicit but data access is complicated. Distributed computing is typically implemented with message passing (e.g. MPI)
Shared Memory:

Data is an a globally accessible address space, any processor can access data by specifying its location using a global index.

Data is mapped out in a natural manner (usually corresponding to the original problem) and access is easy. Information on data locality is obscured and leads to loss of performance.
Global Arrays

Distributed dense arrays that can be accessed through a shared memory-like style

Physically distributed data

single, shared data structure/
global indexing

e.g., access A(4,3) rather than buf(7) on task 2

Global Address Space
Global Arrays (cont.)

- Shared memory model in context of distributed dense arrays
- Much simpler than message-passing for many applications
- Complete environment for parallel code development
- Compatible with MPI
- Data locality control similar to distributed memory/message passing model
- Extensible
- Scalable
Remote Data Access in GA

Message Passing:

identify size and location of data blocks

loop over processors:
  if (me = P_N) then
    pack data in local message buffer
    send block of data to message buffer on P0
  else if (me = P0) then
    receive block of data from P_N in message buffer
    unpack data from message buffer to local buffer
  endif
end loop

copy local data on P0 to local buffer

Global Arrays:

NGA_Get(g_a, lo, hi, buffer, ld);

Global Array handle

Global upper and lower indices of data patch

Local buffer and array of strides
Data Locality

What data does a processor own?

\[
\text{NGA\_Distribution}(g\_a, \text{iproc, lo, hi});
\]

Where is the data?

\[
\text{NGA\_Access}(g\_a, \text{lo, hi, ptr, ld})
\]

Use this information to organize calculation so that maximum use is made of locally held data
Global Array Model of Computations

Shared Object

- copy to local memory
- get

local memory

compute/update

local memory

Shared Object

- copy to shared object
- put

local memory
Example: Matrix Multiply

- local buffers on the processor
- global arrays representing matrices
- `nga_put` and `nga_get`
Matrix Multiply
(a better version)

more scalable!
(less memory, higher parallelism)

atomic accumulate

=  

global

dgemm

local buffers on the processor
Example: 1-D Transpose
Example: 1-D Transpose (cont.)

```c
#define NDIM 1
#define TOTALELEMS 197
#define MAXPROC 128

program main
    implicit none
#include "mafdecls.fh"
#include "global.fh"

integer dims(3), chunk(3), nprocs, me, i, lo(3), hi(3), lo1(3)
integer hi1(3), lo2(3), hi2(3), ld(3), nelem
integer g_a, g_b, a(MAXPROC*TOTALELEMS), b(MAXPROC*TOTALELEMS)
integer heap, stack, ichk, ierr
logical status
heap = 300000
stack = 300000
```
Example: 1-D Transpose (cont.)

c initialize communication library
call mpi_init(ierr)
c initialize ga library
call ga_initialize()
me = ga_nodeid()
nprocs = ga nnodes()
dims(1) = nprocs*TOTALELEMS + nprocs/2 ! Unequal data distribution
ld(1) = MAXPROC*TOTALELEMS
chunk(1) = TOTALELEMS ! Minimum amount of data on each processor
status = ma_init(MT_F_DBL, stack/nprocs, heap/nprocs)

c create a global array
status = nga_create(MT_F_INT, NDIM, dims, "array A", chunk, g_a)
status = ga_duplicate(g_a, g_b, "array B")

c initialize data in GA
do i=1, dims(1)
   a(i) = i
end do
lo1(1) = 1
hi1(1) = dims(1)
if (me.eq.0) call nga_put(g_a, lo1, hi1, a, ld)
call ga_sync() ! Make sure data is distributed before continuing
Example: 1-D Transpose (cont.)

```
c   invert data locally
call nga_distribution(g_a, me, lo, hi)
call nga_get(g_a, lo, hi, a, ld) ! Use locality
nelem = hi(1)-lo(1)+1
do i = 1, nelem
   b(i) = a(nelem - i + 1)
end do

c   invert data globally
lo2(1) = dims(1) - hi(1) + 1
hi2(1) = dims(1) - lo(1) + 1
call nga_put(g_b,lo2,hi2,b,ld)
call ga_sync() ! Make sure inversion is complete
```
Example: 1-D Transpose (cont.)

```plaintext
c check inversion
call nga_get(g_a,lo1,hi1,a,ld)
call nga_get(g_b,lo1,hi1,b,ld)
ichk = 0
do i= 1, dims(1)
   if (a(i).ne.b(dims(1)-i+1).and.me.eq.0) then
      write(6,*), "Mismatch at ",i
      ichk = ichk + 1
   endif
end do
if (ichk.eq.0.and.me.eq.0) write(6,*), "Transpose OK"

status = ga_destroy(g_a) ! Deallocate memory for arrays
status = ga_destroy(g_b)
call ga_terminate()
call mpi_finalize(ierr)
stop
end
```
One-sided Communication

Message Passing:
Message requires cooperation on both sides. The processor sending the message (P1) and the processor receiving the message (P0) must both participate.

One-sided Communication:
Once message is initiated on sending processor (P1) the sending processor can continue computation. Receiving processor (P0) is not involved.
Non-Blocking Communication

- New functionality in GA version 3.3
- Allows overlapping of data transfers and computations
  - Technique for latency hiding
- Nonblocking operations initiate a communication call and then return control to the application immediately
- operation completed locally by making a call to the *wait* routine
**SUMMA Matrix Multiplication**

\[
C = A \cdot B
\]

**Computation**

- Issue NB Get A and B blocks
- \( \text{do} \) (until last chunk)
  - issue NB Get to the next blocks
  - wait for previous issued call
  - compute \( A \cdot B \) (sequential dgemm)
- NB atomic accumulate into “C” matrix
- \( \text{done} \)

**Advantages:**
- Minimum memory
- Highly parallel
- Overlaps computation and communication
- Latency hiding
- Exploits data locality
- Patch matrix multiplication (easy to use)
- Dynamic load balancing
SUMMA Matrix Multiplication: Improvement over MPI

Non-Blocking Communication Performance

*2.4Ghz P4 Linux cluster, Myrinet-GM interconnect (at SUNY, Buffalo)*
Structure of GA

Application programming language interface

Global Arrays and MPI are completely interoperable. Code can contain calls to both libraries.

- Fortran 77
- C
- C++
- Python
- Babel

Global Arrays and MPI are completely interoperable. Code can contain calls to both libraries.

- Message Passing
  - Global operations

- ARMCI
  - portable 1-sided communication
  - put, get, locks, etc

- System specific interfaces
  - LAPI, GM/Myrinet, threads, VIA, ...
Core Capabilities

- **Distributed array library**
  - dense arrays 1-7 dimensions
  - four data types: `integer`, `real`, `double precision`, `double complex`
  - global rather than per-task view of data structures
  - user control over data distribution: regular and irregular

- **Collective and shared-memory style operations**
  - `ga_sync`, `ga_scale`, etc
  - `ga_put`, `ga_get`, `ga_acc`
  - nonblocking `ga_put`, `ga_get`, `ga_acc`

- **Interfaces to third party parallel numerical libraries**
  - PeIGS, Scalapack, SUMMA, Tao
    - example: to solve a linear system using LU factorization
    ```
call ga_lu_solve(g_a, g_b)
    ```
    instead of
    ```
call pdgetrf(n,m, locA, p, q, dA, ind, info)
call pdgetrs(trans, n, mb, locA, p, q, dA, dB, info)
    ```
Interoperability and Interfaces

- Language interfaces to Fortran, C, C++, Python
- Interoperability with MPI and MPI libraries
  - e.g., PETSC, CUMULVS
- Explicit interfaces to other systems that expand functionality of GA
  - ScaLAPACK-scalable linear algebra software
  - Peigs-parallel eigensolvers
  - TAO-advanced optimization package
Global Array Processor Groups

Many parallel applications require the execution of a large number of independent tasks. Examples include

- Numerical evaluation of gradients
- Monte Carlo sampling over initial conditions or uncertain parameter sets
- Free energy perturbation calculations (chemistry)
- Nudged elastic band calculations (chemistry and materials science)
- Sparse matrix-vector operations (NAS CG benchmark)
Global Array Processor Groups

If the individual calculations are small enough then each processor can be used to execute one of the tasks (embarrassingly parallel algorithms).

If the individual tasks are large enough that they must be distributed amongst several processors then the only option (usually) is to run each task sequentially on multiple processors. This usually limits the total number of processors that can be applied to the problem since parallel efficiency degrades as the number of processors increases.
Alternatively the collection of processors can be decomposed into processor groups. These processor groups can be used to execute parallel algorithms *independently* of one another. This requires

- global operations that are restricted in scope to a particular group instead of over the entire domain of processors (world group)

- distributed data structures that are restricted to a particular group
integer function ga_pgroup_create(list, count)

Returns a handle to a group of processors. The total number of processors is count, the individual processor IDs are located in the array list.

subroutine ga_pgroup_set_default(p_grp)

Set the default processor to p_grp. All arrays created after this point are created on the default processor group, all global operations are restricted to the default processor group unless explicit directives are used. Initial value of the default processor group is the world group.
Explicit Operations on Groups

Explicit Global Operations on Groups

\begin{align*}
ga\_{\text{pgroup}}\_\text{sync}(p\_\text{grp}) \\
ga\_{\text{pgroup}}\_\text{brdcst}(p\_\text{grp}, \text{type}, \text{buf}, \text{lenbuf}, \text{root}) \\
ga\_{\text{pgroup}}\_\text{igop}(p\_\text{grp}, \text{type}, \text{buf}, \text{lenbuf}, \text{op}) \\
ga\_{\text{pgroup}}\_\text{dgop}(p\_\text{grp}, \text{type}, \text{buf}, \text{lenbuf}, \text{op})
\end{align*}

Query Operations on Groups

\begin{align*}
ga\_{\text{pgroup}}\_\text{nnodes}(p\_\text{grp}) \\
ga\_{\text{pgroup}}\_\text{nodeid}(p\_\text{grp})
\end{align*}

Access Functions

\begin{align*}
\text{integer function } ga\_{\text{pgroup}}\_\text{get}\_\text{default}() \\
\text{integer function } ga\_{\text{pgroup}}\_\text{get}\_\text{world}()
\end{align*}
Programming with Groups

- Most explicit group operations in GA reflect operations available for MPI groups
- Concept of default group is not available in MPI
- Higher level abstractions not available in MPI
Copy and copy_patch operations are supported for global arrays that are created on different groups. One of the groups must be completely contained in the other (nested).

The copy or copy_patch operation must be executed by all processors on the nested group (group B in illustration)
c set up groups
me = ga_nodeid()
nprocs = ga_nnodes()
grpsize = 4
ngrps = nprocs/grpsize
nproc = grpsize
do i = 1, ngrps  ! All processors participate in
  do j = 1, grpsize ! creation of group
    proclist(j) = grpsize*(i-1) + (j-1)
  end do
  procgroup(i) = ga_pgroup_create(proclist,nproc)
end do
my_pgrp = (me - mod(me,grpsize))/grpsize + 1

! run task on groups
!call ga_pgroup_set_default(procgroup(my_pgrp))
call do_parallel_task
call ga_pgroup_set_default(ga_pgroup_get_world())
MD Example

Spatial Decomposition Algorithm:

- Partition particles among processors
- Update coordinates at every step
- Update partitioning after fixed number of steps
MD Parallel Scaling

Scaling of Single Parallel Task

- Single Parallel Task
- Perfect Scaling

Number of Processors

Speedup
MD Performance on Groups

Scaling of Parallel MD Tasks on Groups

- 256 Tasks
- 1024 Tasks
- Perfect Scaling

Number of Processors vs. Speedup
Application Areas

bioinformatics
electronic structure chemistry
  GA is the standard programming model
glass flow simulation
biology

thermal flow simulation
material sciences
molecular dynamics
Visualization and image analysis

Others: financial security forecasting, astrophysics, geosciences, atmospheric chemistry
Ghost Cells

normal global array

global array with ghost cells

Operations:

- **NGA_Create_ghosts**: creates array with ghosts cells
- **GA_Update_ghosts**: updates with data from adjacent processors
- **NGA_Access_ghosts**: provides access to “local” ghost cell elements
- **NGA_Nbget_ghost_dir**: nonblocking call to update ghosts cells
Ghost Cell Update

Automatically update ghost cells with appropriate data from neighboring processors. A multiprotocol implementation has been used to optimize the update operation to match platform characteristics.
\[ f_i(r + e_i, t + \Delta t) = f_i(r, t) - \frac{1}{\tau} (f_i(r, t) - f_i^{eq}(r, t)) \]
Ghost Cell Application Performance

![Graph showing the relationship between time and number of processors for Total and Update activities.](image)

**Graph:**
- **Y-axis:** Time (sec)
- **X-axis:** Number of Processors
- **Legend:**
  - Red dots: Total
  - Blue squares: Update

**Image:**
- Visualization of a simulation or data visualization related to ghost cell application performance.
Mirrored Arrays

- Create Global Arrays that are replicated between SMP nodes but distributed within SMP nodes
- Aimed at fast nodes connected by relatively slow networks (e.g. Beowulf clusters)
- Use memory to hide latency
- Most of the operations supported on ordinary Global Arrays are also supported for mirrored arrays
- Global Array toolkit augmented by a merge operation that adds all copies of mirrored arrays together
- Easy conversion between mirrored and distributed arrays
Mirrored Arrays (cont.)
NWChem DFT Calculation

http://www.emsl.pnl.gov/docs/nwchem
Other Functionality

- Common Component Architecture
- Disk Resident Arrays
  - Provide an interface between GA and distributed files on disk
- Sparse data manipulation
Related Programming Tools

- Co-Array Fortran
  - Distributed Arrays
  - One-Sided Communication
  - No Global View of Data

- UPC
  - Model Similar to GA but only applicable to C programs
  - Global Shared Pointers could be used to implement GA functionality
    - C does not really support multi-dimensional arrays

- High level functionality in GA is missing from these systems
Summary

· The idea has proven very successful
  · efficient on a wide range of architectures
    · core operations tuned for high performance
  · library substantially extended but all original (1994) APIs preserved
  · increasing number of application areas

· Supported and portable tool that works in real applications

· Future work
  · Fault tolerance
Source Code and More Information

- Version 3.3 available
- Version 3.4 (with groups) available in beta
- Homepage at http://www.emsl.pnl.gov/docs/global/
- Platforms (32 and 64 bit)
  - IBM SP
  - Cray X1, XD1
  - Linux Cluster with Ethernet, Myrinet, Infiniband, or Quadrics
  - Solaris
  - Fujitsu
  - Hitachi
  - NEC
  - HP
  - Windows
Disk Resident Arrays

- Extend GA model to disk
  - similar system to Panda (U. Illinois) but higher level APIs
- Provide easy transfer of data between N-dim arrays stored on disk and distributed arrays stored in memory
- Use when
  - Arrays too big to store in core
  - checkpoint/restart
  - out-of-core solvers
High Bandwidth Read/Write

Disk Resident Array

Disk Resident Arrays automatically decomposed into multiple files

Disks
Scalable Performance of DRA

SMP node

I/O buffers

file systems

bar chart showing bandwidth [MB/s] vs. array rank and disks.
Useful GA Functions (Fortran)

```fortran
subroutine ga_initialize()
subroutine ga_terminate()

integer function ga_nnodes()
integer function ga_nodeid()

logical function nga_create(type, dim, dims, name, chunk, g_a)
    integer type (MT_F_INT, MT_F_DBL, etc.)
    integer dim
    integer dims(dim)
    character(*) name
    integer chunk(dim)
    integer g_a

logical function ga_duplicate(g_a, g_b, name)
    integer g_a
    integer g_b
    character(*) name

logical function ga_destroy(g_a)
    integer g_a

subroutine ga_sync()
```
Use GA Functions (Fortran)

```fortran
subroutine nga_distribution(g_a, node_id, lo, hi)
  integer g_a
  integer node_id
  integer lo(dim)
  integer hi(dim)
end subroutine nga_distribution

subroutine nga_put(g_a, lo, hi, buf, ld)
  integer g_a
  integer lo(dim)
  integer hi(dim)
  fortran array buf
  integer ld(dim-1)
end subroutine nga_put

subroutine nga_get(g_a, lo, hi, buf, ld)
  integer g_a
  integer lo(dim)
  integer hi(dim)
  fortran array buf
  integer ld(dim-1)
end subroutine nga_get
```
void GA_Initialize()
void GA_Terminate()

int GA_Nnodes()
int GA_Nodeid()

int NGA_Create(type, dim, dims, name, chunk)
    int type (C_INT, C_DBL, etc.)
    int dim
    int dims[dim]
    char* name
    int chunk[dim]
    Returns GA handle g_a
int GA_Duplicate(g_a, name)
    int g_a
    Returns GA handle g_b
    char* name
void GA_Destroy(g_a)
    int g_a

void GA_Sync()}
void NGA_Distribution(g_a, node_id, lo, hi)
    int g_a
    int node_id
    int lo[dimension]
    int hi[dimension]

void NGA_Put(g_a, lo, hi, buf, ld)
    int g_a
    int lo[dimension]
    int hi[dimension]
    void* buf
    int ld[dimension-1]

void NGA_Get(g_a, lo, hi, buf, ld)
    int g_a
    int lo[dimension]
    int hi[dimension]
    void* buf
    int ld[dimension-1]