Flexible Hierarchical Execution of Parallel Task Loops

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Injection Bandwidth vs CPU speeds





Kale (Salishan 2018)



Motivation

- Trend:
 - Deeper nodes
 - Thinner pipes
- Accelerators (e.g. GPUs)
- Increased Programmer effort

Year	Machine	Linpack (FLOPs)	FLOPs/ Local	FLOPs/ Remote
1988	Cray YMP	2.1 Giga	0.52	0.52
1997	ASCI Red	1.6 Tera	8.3	20
2011	Road- runner	1.0 Peta	6.7	170
2012	Sequoia	17 Peta	32	160
2013	Titan	18 Peta	29	490
2018	Summit	122 Peta	37	1060
2011	K-Comp	11 Peta	15	95
2013	Tianhe-2	34 Peta	22	1500
2016	Sunway	93 Peta	130	1500
2021	TBD	1.0 Exa	80	3200
2021	TBD	1.0 Exa	300	10000





Fat Nodes

First law of holes:

• If you find yourself in a hole, stop digging!







Main Idea: Spreading Work Across Cores

- Speed up individual calculations via OpenMP
- FLOPs are cheap, need to inject early
- Better communication, computation overlap









Motivation

New Axes of Optimization

- Problem Size Decomposition (Grain Size)
- Resources Assigned to a Task (Spreading)





Experimental Setup

- Charm Build
 - Separate processes (Non-SMP mode)
 - -O3 -with-production
 - PAMI-LRTS communication layer
- Five Runs
 - OpenMP Threads (Spreading) = 1, 2, etc
 - Grid Size = 178848² doubles (~90%)
 - Block Size = 7452, various
 - Chares (Objects) = 24²
 - Iterations = 10-100
 - Nodes = 4





OpenMP Pragmas

- Schedule Static
- Chunk Size (Iterations)
 - Default (Block / Cores)
 - 1
 - 16
 - 512
- Collapse





Machines

Bridges (PSC)

- 2 x 14-core Haswell E5-2695
- 128 GB DDR4

Summit (ORNL)

- 2 x 22-core IBM Power9
- 512 GB DDR4













What happens when we eliminate communication?

i.e. are effects just from improved caching?







Lets look at communication performance...

using projections.





OpenMP Baseline



Charm++ Baseline



Time (s)

Spreading Technique









Runtime Integration





Automating teams configuration

• Broader Agenda

- Automate decisions -> easier for user
- "Spread": How many teams, i.e how many masters and how many drones?
- Other runtime decisions:
- How many ppn, i.e cores per process?
- How many processes per node?
- How many cores to turn off (memory bottleneck)?
- Enable SMT or not?





Automating teams configuration

- Use OpenMP to create master thread on all cores
- Integrate with load balancing framework to change master thread count
- Use OpenMP nested parallelism to set/change number of drone threads within the application
 - Use pthread affinities instead of OpenMP affinity to update configurations at runtime
- Runtime selects the best performing configuration after testing with different configurations (one per LB step)





Using OpenMP with nested parallelism (static)

Bridges - single-node integrated OpenMP runs for SMP and Non-SMP builds







Using OpenMP with nested parallelism (static)

Stampede2 - Skylake 4-node run integrated OpenMP







OpenMP Implementation

```
machine-smp.C
```

```
int num_threads = tocreate + 1;
omp_set_dynamic(0);
omp_set_num_threads(num_threads);
#pragma omp parallel
```

```
{
   size_t i = omp_get_thread_num();
   call_startfn((void *)i);
}
```

```
jacobi2d.C
```

Static configuration: OMP_NESTED=true GOMP_CPU_AFFINITY=0-7 (eg. for 8 cores) OMP_PROC_BIND=spread

Dynamic configuration: pthread_setaffinity_np(thread, sizeof(cpu_set_t), &cpuset);





OpenMP implementation with pthread affinity

- Similar performance with processbased and OpenMP implementations
 - Some NUMA effects
- OpenMP Limitations:
 - Nested parallelism configurations cannot be dynamically changed
 - Affinities are set at the initialization and cannot be changed
- With Charm++ we are able to dynamically change OpenMP configurations and with pthread affinity we set affinities for each new configuration







Next steps

- Integrate the LB framework to fully automate configuration selection
 - Current implementation is able to dynamically set different configurations at runtime based on user input
 - Benefit over static OpenMP configuration configurations and affinities can be changed at runtime
- Compare with CkLoop implementation in Charm++





Summary

- Spreading offers new optimization parameter
- Increases performance 20-30% in prototype application
- Spread factor is controllable at runtime
- Current integration into Charm++ ongoing
 Questions

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