Advanced Charm++ Tutorial

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Topics For This Talk

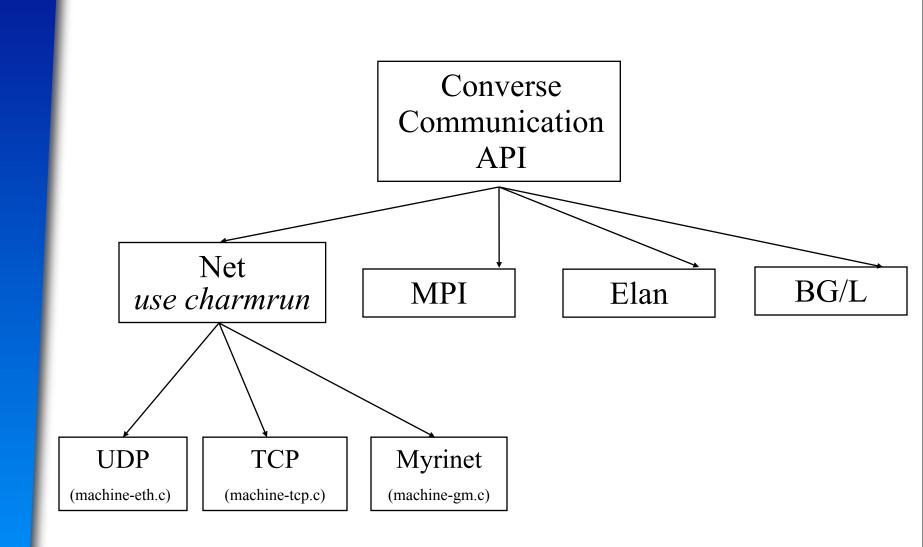
- Building Charm++
- Advanced messaging
- Interface file (.ci)
- Advanced load balancing
- Groups
- Threads
- Delegation
- Array multicast
- SDAG

Charm++ on Parallel Machines

Runs on:

- Any machine with MPI, including
 - IBM Blue Gene/L, SP
 - Cray XT3
 - SGI Altix
- PSC's Lemieux (Quadrics Elan)
- Clusters with Ethernet (UDP/TCP)
- Clusters with Myrinet (GM or MX)
- Apple clusters
- Even Windows!
- SMP-Aware (pthreads)

Communication Architecture



./build

Usage: build **<target>** <version> <options> [charmc-options ...]

<target>: converse charm++ LIBS AMPI FEM bigemulator pose jade msa doc ps-doc pdf-doc html-doc

charm++ compile Charm++ core only

AMPI compile Adaptive MPI on top of Charm++

FEM compile FEM framework

LIBS compile additional parallel libraries with Charm++ core

bigemulator build additional BigSim libraries

pose build POSE parallel discrete event simulator

jade build Jade compiler (auto-builds charm++, msa)

msa build Multiphase Shared Arrays(MSA) library

./build

Usage: build <target> **<version>** <options> [charmc-options ...]

<version>: Basic configurations

bluegenel mpi-sp net-sol-x86

elan-axp ncube2 net-sun elan-linux-ia64 net-axp net-win32

exemplar net-cygwin net-win64

mpi-axp net-darwin-x86 origin-pthreads

mpi-bluegenel net-hp origin2000

mpi-crayx1 net-hp-ia64 portals-crayxt3

mpi-crayxt3 net-irix shmem-axp

mpi-exemplar net-linux sim-linux

mpi-hp-ia64 net-linux-amd64 sp3
mpi-linux net-linux-axp t3e

mpi-linux-amd64 net-linux-cell uth-linux

mpi-linux-axp net-linux-ia64 uth-win32

mpi-linux-ia64 net-linux-ppc vmi-linux

mpi-origin net-ppc-darwin vmi-linux-amd64
mpi-ppc-darwin net-rs6k vmi-linux-ia64

mpi-ppc-darwin net-rs6k vmi-linux-ia64
mpi-sol net-sol

mpi-sol-amd64 net-sol-amd64

./build

Usage: build <target> **<version>** <options> [charmc-options ...]

<version>: Basic configurations

bluegenel

elan-axp

elan-linux-ia64

exemplar

mpi-axp

mpi-bluegenel

mpi-crayx1

mpi-crayxt3

mpi-exemplar

mpi-hp-ia64

mpi-linux

mpi-linux-amd64

mpi-linux-axp

mpi-linux-ia64

mpi-origin

mpi-ppc-darwin

mpi-sol

mpi-sol-amd64

mpi-sp

ncube2

net-axp

net-cygwin

net-darwin-x86

net-hp

net-hp-ia64

net-irix

net-linux

net-linux-amd64

net-linux-axp

net-linux-cell

net-linux-ia64

net-linux-ppc

net-ppc-darwin

net-rs6k

net-sol

net-sol-amd64

net-sol-x86

net-sun

net-win32

net-win64

origin-pthreads

origin2000

portals-crayxt3

shmem-axp
sim-linux

sp3

t3e

uth-linux

uth-win32

vmi-linux

vmi-linux-amd64

vmi-linux-ia64

./build

Usage: build <target> <version> **<options>** [charmc-options ...]

<options>: compiler and platform specific options

Platform specific options (choose multiple if they apply):

lam	Use LAM MPI
smp	support for SMP, multithreaded charm on each node
mpt	use SGI Message Passing Toolkit (only for mpi version)
gm	use Myrinet for communication
tcp	use TCP sockets for communication (ony for net version)
vmi	use NCSA's VMI for communication (only for mpi version)
scyld	compile for Scyld Beowulf cluster based on bproc
clustermatic	compile for Clustermatic (support version 3 and 4)
pthreads	compile with pthreads Converse threads

./build

Usage: build <target> <version> **<options>** [charmc-options ...]

<options>: compiler and platform specific options

Advanced options:

bigemulator compile for BigSim simulator

ooc compile with out of core support

syncft compile with Charm++ fault tolerance support

papi compile with PAPI performance counter support (if any)

Charm++ dynamic libraries:

--build-shared build Charm++ dynamic libraries (.so) (default)

--no-build-shared don't build Charm++'s shared libraries

```
./build
```

Usage: build <target> <version> **<options>** [charmc-options ...]

<options>: compiler and platform specific options

Choose a C++ compiler (only one option is allowed from this section):

cc, cc64 For Sun WorkShop C++ 32/64 bit compilers

cxx DIGITAL C++ compiler (DEC Alpha)

kcc KAI C++ compiler

pgcc Portland Group's C++ compiler

acc HP aCC compiler

icc Intel C/C++ compiler for Linux IA32

ecc Intel C/C++ compiler for Linux IA64

gcc3 use gcc3 - GNU GCC/G++ version 3

gcc4 use gcc4 - GNU GCC/G++ version 4 (only mpi-crayxt3)

mpcc SUN Solaris C++ compiler for MPI

pathscale use pathscale compiler suite

./build

Usage: build <target> <version> **<options>** [charmc-options ...]

<options>: compiler and platform specific options

Choose a fortran compiler (only one option is allowed from this section):

g95 G95 at http://www.g95.org

absoft Absoft fortran compiler

pgf90 Portland Group's Fortran compiler

ifc Intel Fortran compiler (older versions)

ifort Intel Fortran compiler (newer versions)

./build

Usage: build <target> <version> <options> [charmc-options ...]

<charmc-options>: normal compiler options
-g -0 -save -verbose

To see the latest versions of these lists or to get more detailed help, run ./build --help

Build Script

Build script does:

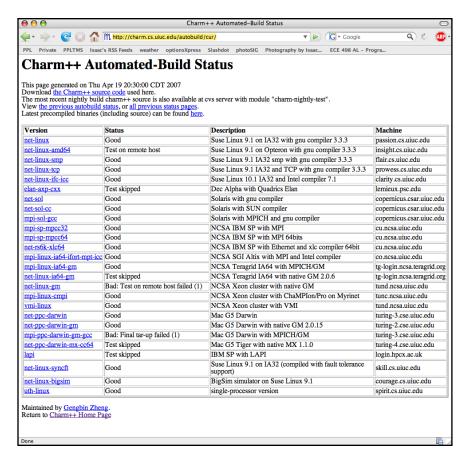
- ./build <target> <version> <options> [charmc-options ...]
- Creates directories <version> and <version>/tmp
- Copies src/scripts/Makefile into <version>/tmp
- Does a
 "make <target> <version> OPTS=<charmc-options>"
 in <version>/tmp
- That's all build does. The rest is handled by the Makefile.

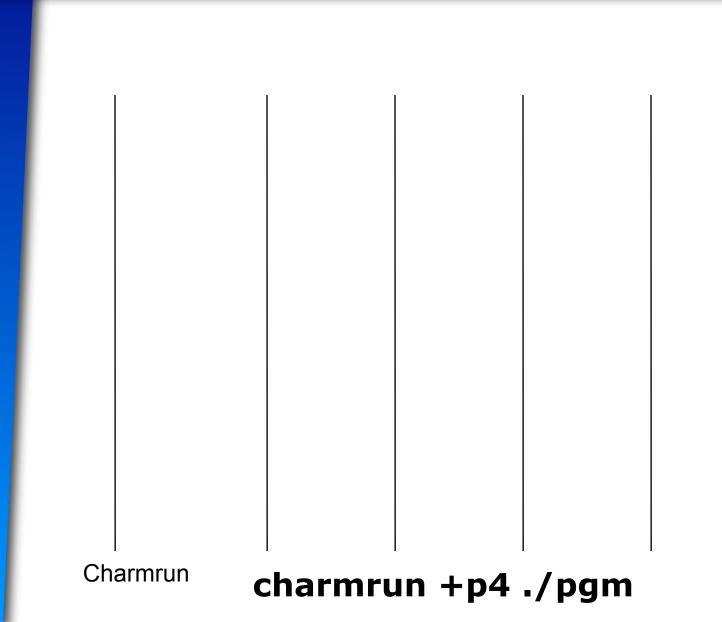
How 'build' works

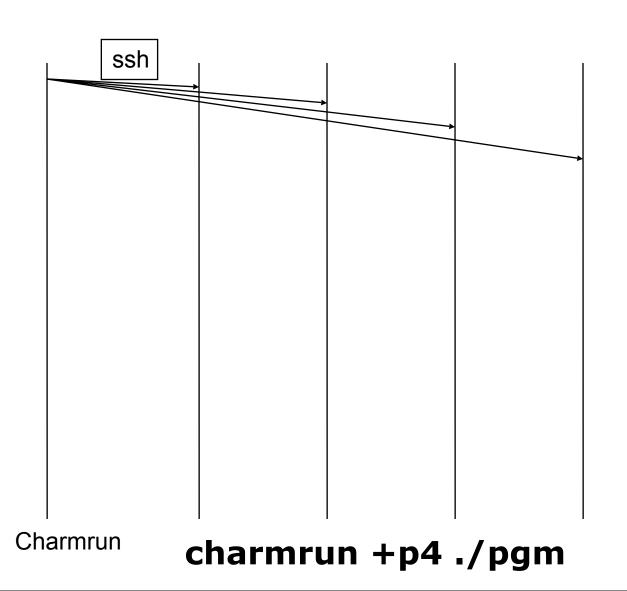
- build AMPI net-linux gm kcc
 - Mkdir net-linux-gm-kcc
 - Cat conv-mach-[kcc|gm|smp].h to conv-mach-opt.h
 - Cat conv-mach-[kcc|gm].sh to conv-mach-opt.sh
 - Gather files from net, etc (Makefile)
 - Make charm++ under
 - net-linux-gm/tmp

What if build fails?

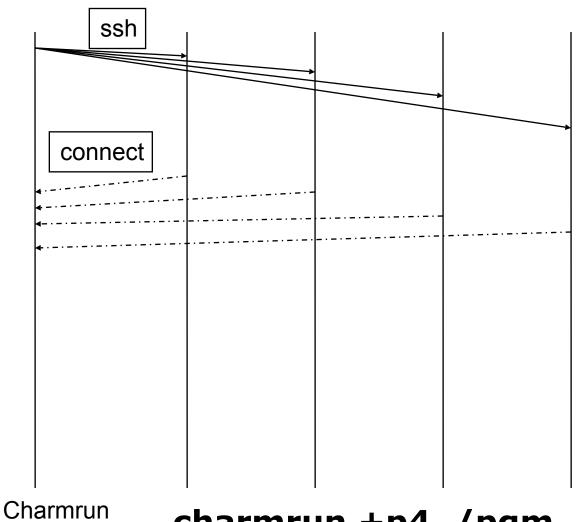
- Use latest version from CVS
- Check the nightly auto-build tests: http://charm.cs.uiuc.edu/autobuild/cur/
- Email: ppl@cs.uiuc.edu



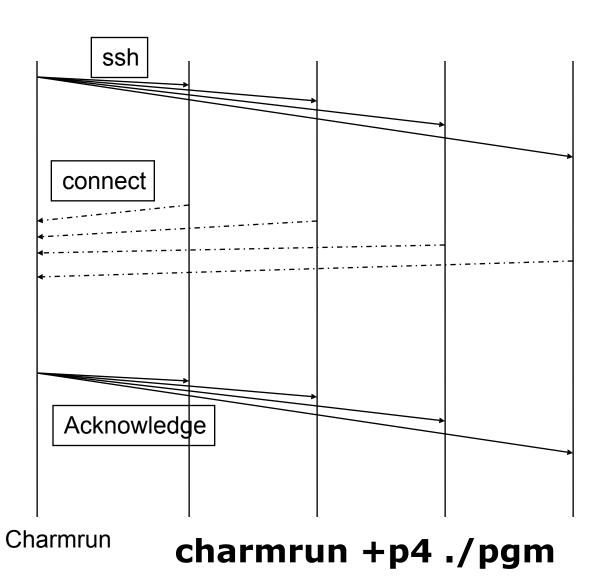




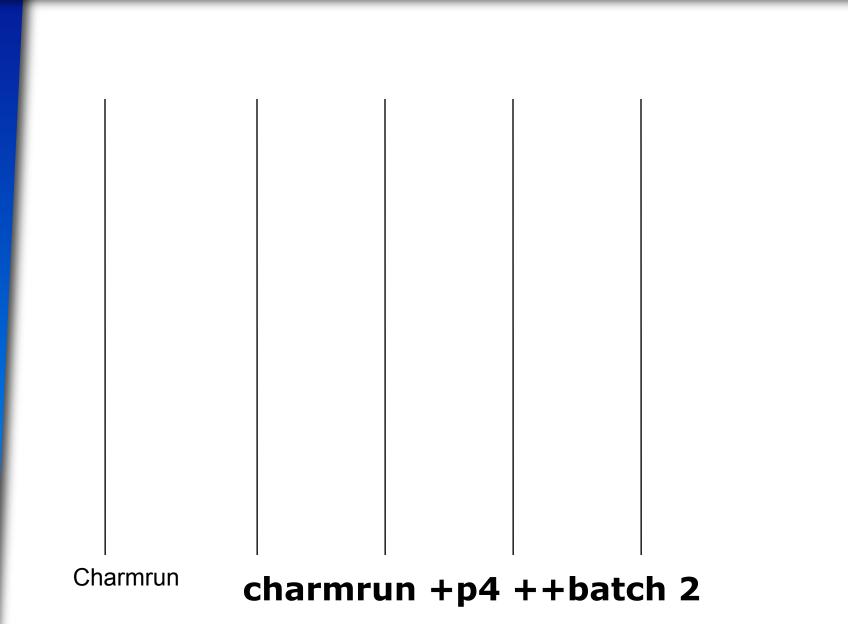
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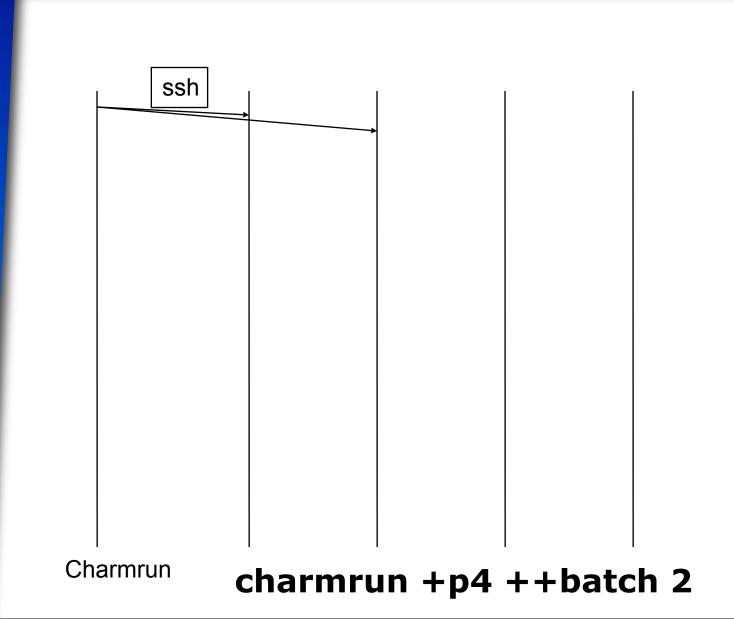


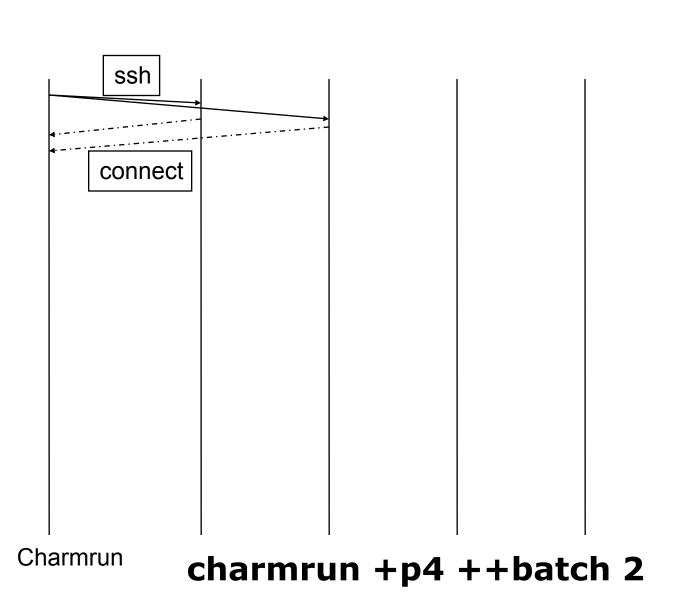
charmrun +p4 ./pgm

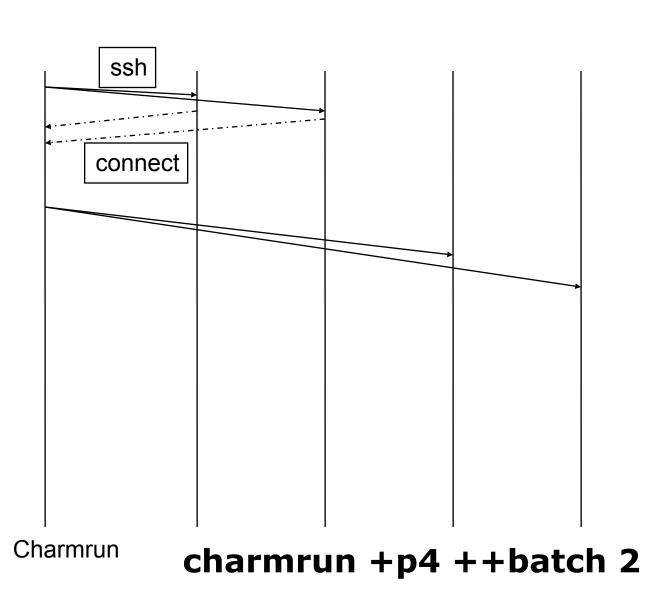


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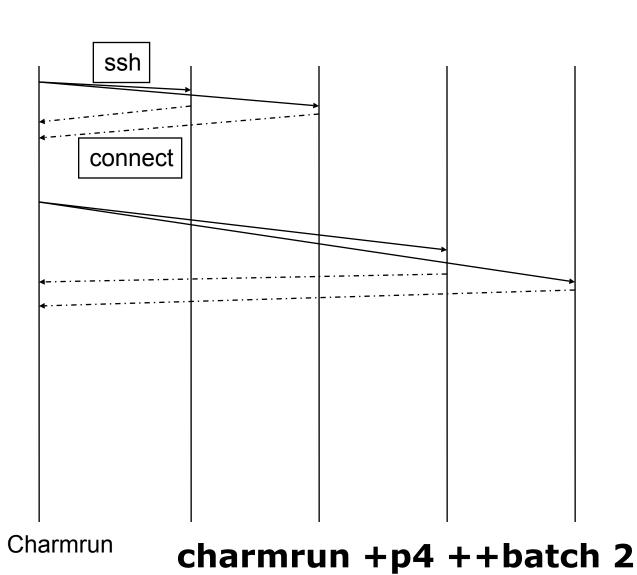




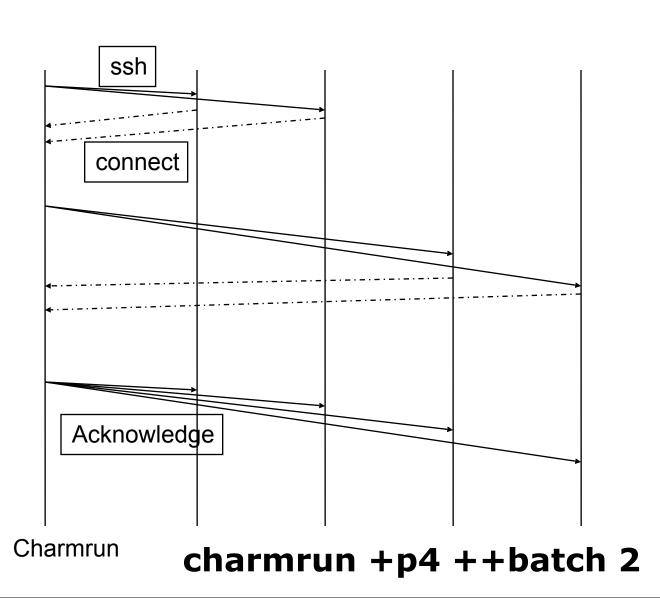




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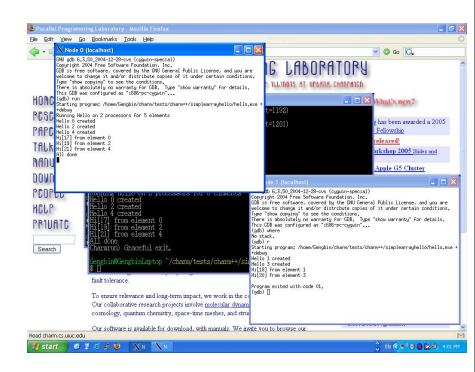


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Debugging Charm++ Applications

- printf
- Gdb
 - Sequentially (standalone mode)
 - gdb ./pgm +vp16
 - Attach gdb manually
 - Run debugger in xterm
 - charmrun +p4 pgm ++debug
 - charmrun +p4 pgm ++debug-no-pause
 - Memory paranoid
 - -memory paranoid
 - Parallel debugger



Advanced Messaging



Prioritized Execution

- Charm++ scheduler
 - Default FIFO (oldest message)
- Prioritized execution
 - If several messages available, Charm will process the messages in the order of their priorities
- Very useful for speculative work, ordering timestamps, etc...

Priority Classes

- Charm++ scheduler has three queues: high, default, and low
- As signed integer priorities:
 - High -MAXINT to -1
 - Default 0
 - Low 1 to +MAXINT
- As unsigned bitvector priorities:
 - 0x0000 Highest priority -- 0x7FFF
 - 0x8000 Default priority
 - 0x8001 -- 0xFFFF Lowest priority

Prioritized Messages

Number of priority bits passed during message allocation

```
FooMsg * msg = new (size, nbits) FooMsg;
```

- Priorities stored at the end of messages
- Signed integer priorities

```
*CkPriorityPtr(msg)=-1;
CkSetQueueing(msg, CK_QUEUEING_IFIFO);
```

Unsigned bitvector priorities

```
CkPriorityPtr(msg)[0]=0x7fffffff;
CkSetQueueing(msg, CK_QUEUEING_BFIF0);
```

Prioritized Marshalled Messages

- Pass "CkEntryOptions" as last parameter
- For signed integer priorities:

```
CkEntryOptions opts;
opts.setPriority(-1);
fooProxy.bar(x,y,opts);
```

For bitvector priorities:

```
CkEntryOptions opts;
unsigned int prio[2]={0x7FFFFFFF,0xFFFFFFF};
opts.setPriority(64,prio);
fooProxy.bar(x,y,opts);
```

Advanced Message Features

- Read-only messages
 - Entry method agrees not to modify or delete the message
 - Avoids message copy for broadcasts, saving time
- Inline messages
 - Direct method invocation if on local processor
- Expedited messages
 - Message do not go through the charm++ scheduler (ignore any Charm++ priorities)
- Immediate messages
 - Entries are executed in an interrupt or the communication thread
 - Very fast, but tough to get right
 - Immediate messages only currently work for NodeGroups and Group (non-smp)
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Read-Only, Expedited, Immediate

All declared in the .ci file

```
entry [nokeep] void foo_readonly(Msg *);
entry [inline] void foo_inl(Msg *);
entry [expedited] void foo_exp(Msg *);
entry [immediate] void foo_imm(Msg *);
...
};
```

Interface File (ci)



Interface File Example

```
mainmodule hello {
  include "myType.h"
  initnode void myNodeInit();
  initproc void myInit();
  mainchare mymain {
    entry mymain(CkArgMsg *m);
  };
  array[1D] foo {
    entry foo(int problemNo);
    entry void bar1(int x);
    entry void bar2(myType x);
  };
};
```

Include and Initcall

- Include
 - Include an external header files
- Initcall
 - User plugging code to be invoked in Charm++'s startup phase
 - Initnode
 - Called once on every node
 - Initproc
 - Called once on every processor
 - Initnode calls are called before Initproc calls

Entry Attributes

- Threaded
 - Function is invoked in a CthThread
- Sync
 - Blocking methods, can return values as a message
 - Caller must be a thread
- Exclusive
 - For Node Group
 - Do not execute while other exclusive entry methods of its node group are executing in the same node
- Notrace
 - Invisible to trace projections
 - entry [notrace] void recvMsg(multicastGrpMsg *m);

Groups/Node Groups



Groups and Node Groups

- Groups
 - Similar to arrays:
 - Broadcasts, reductions, indexing
 - But not completely like arrays:
 - Non-migratable; one per processor
 - Exactly one representative on each processor
 - Ideally suited for system libraries
 - Historically called branch office chares (BOC)
- Node Groups
 - One per SMP node

Declarations

.ci file

```
group mygroup {
    entry mygroup(); //Constructor
    entry void foo(foomsg *); //Entry method
  };
  nodegroup mynodegroup {
     entry mynodegroup(); //Constructor
     entry void foo(foomsg *); //Entry method
 };
C++ file
 class mygroup : public Group {
            mygroup() {}
    void foo(foomsg *m) { CkPrintf("Do Nothing");}
 };
 class mynodegroup : public NodeGroup {
            mynodegroup() {}
    void foo(foomsg *m) { CkPrintf("Do Nothing");}
 };
```

Creating and Calling Groups

Creation

```
p = CProxy mygroup::ckNew();
```

Remote invocation

```
p.foo(msg); //broadcast
p[1].foo(msg); //asynchronous
p.foo(msg, npes, pes); // list send
```

Direct local access

```
mygroup *g=p.ckLocalBranch();
g->foo(....); //local invocation
```

Danger: if you migrate, the group stays behind!

Advanced Load-balancers



Advanced load balancing: Writing a new strategy

Inherit from CentralLB and implement the work(...) function

LB Database

```
struct LDStats {
   ProcStats *procs;
   LDObjData* objData;
   LDCommData* commData;
   int *to_proc;
   //.. .. ..
//Dummy Work function which assigns all objects to
//processor 0
//Don't implement it!
void fooLB::work(CentralLB::LDStats* stats,int count){
     for(int count=0;count < nobjs; count++)</pre>
        stats.to proc[count] = 0;
```

Compiling and Integration

- Edit and run Makefile_lb.sh
 - Creates Make.lb which is included by the main Makefile
- Run make depends to correct dependencies
- Rebuild charm++ and is now available in -balancer fooLB

Threads in Charm++



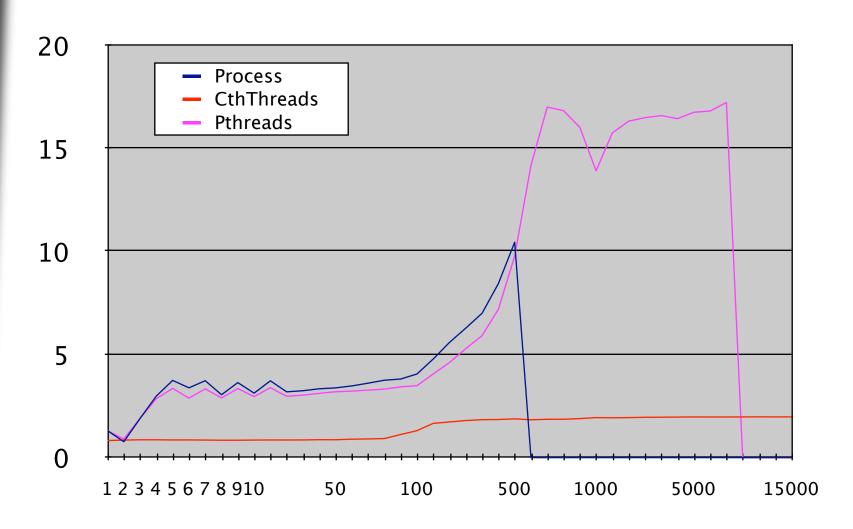
Why use Threads?

- They provide one key feature: blocking
 - Suspend execution (e.g., at message receive)
 - Do something else
 - Resume later (e.g., after message arrives)
- Example: MPI_Recv, MPI_Wait semantics
- Function call interface more convenient than message-passing
 - Regular call/return structure (no CkCallbacks) with complete control flow
 - Allows blocking in middle of deeply

Why <u>not</u> use Threads?

- Slower
 - Around 1us context-switching overhead unavoidable
 - Creation/deletion perhaps 10us
- Migration more difficult
 - State of thread is scattered through stack, which is maintained by compiler
 - By contrast, state of object is maintained by users
- Thread disadvantages form the motivation to use SDAG (later)

Context Switch Cost



What are (Converse) Threads?

- One flow of control (instruction stream)
 - Machine Registers & program counter
 - Execution stack
- Like pthreads (kernel threads)
- Only different:
 - Implemented at user level (in Converse)
 - Scheduled at user level; non-preemptive
 - Migratable between nodes

How do I use Threads?

Many options:

- AMPI
 - Always uses threads via TCharm library
- Charm++
 - [threaded] entry methods run in a thread
 - [sync] methods
- Converse
 - C routines CthCreate/CthSuspend/CthAwaken
 - Everything else is built on these
 - Implemented using
 - SYSV makecontext/setcontext
 - POSIX setjmp/alloca/longjmp
 - Assembly code

How do I use Threads (example)

Blocking API routine: find array element

```
int requestFoo(int src) {
  myObject *obj=...;
  return obj->fooRequest(src)
}
```

Send request and suspend

```
int myObject::fooRequest(int src) {
  proxy[dest].fooNetworkRequest(thisIndex);
  stashed_thread=CthSelf();
  CthSuspend(); // -- blocks until awaken call --
  return stashed_return;
}
```

Awaken thread when data arrives

```
void myObject::fooNetworkResponse(int ret) {
   stashed_return=ret;
   CthAwaken(stashed_thread);
}
```

How do I use Threads (example)

Send request, suspend, recv, awaken, return

```
int myObject::fooRequest(int src) {
  proxy[dest].fooNetworkRequest(thisIndex);
  stashed_thread=CthSelf();
  CthSuspend();
     void myObject::fooNetworkResponse(int ret) {
        stashed_return=ret;
        CthAwaken(stashed_thread);
    }
  return stashed return;
```

Thread Migration



Stack Data

- The stack is used by the compiler to track function calls and provide temporary storage
 - Local Variables
 - Subroutine Parameters
 - C "alloca" storage
- Most of the variables in a typical application are stack data
- Stack is allocated by Charm run-time as heap memory (+stacksize)

Migrate Stack Data

- Without compiler support, cannot change stack's address
 - Because we can't change stack's interior pointers (return frame pointer, function arguments, etc.)
- Existing pointers to addresses in original stack become invalid
- Solution: "isomalloc" addresses
 - Reserve address space on every processor for every thread stack
 - Use mmap to scatter stacks in virtual memory efficiently
 - Idea comes from PM²

Migrate Stack Data

Processor A's Memory

Processor B's Memory

0xFFFFFFF

Thread 1 stack

0xFFFFFFF

Thread 2 stack

Thread 3 stack

Thread 4 stack

Heap

Globals

Code

<u>Migrate</u> Thread 3

Heap

Globals

Code

0x00000000

0x0000000

49

Migrate Stack Data: Isomalloc

<u>Migrate</u>

Thread 3

Processor A's Memory

Processor B's Memory

0xFFFFFFF

Thread 2 stack

0xFFFFFFF

Thread 4 stack

Heap

Globals

Code

Thread 1 stack

Thread 3 stack

Heap

Globals

Code

0x00000000

0x00000000

50

Migrate Stack Data

- Isomalloc is a completely automatic solution
 - No changes needed in application or compilers
 - Just like a software shared-memory system, but with proactive paging
- But has a few limitations
 - Depends on having large quantities of virtual address space (best on 64-bit)
 - 32-bit machines can only have a few gigs of isomalloc stacks across the whole machine
 - Depends on unportable mmap
 - Which addresses are safe? (We must guess!)
 - What about Windows? Or Blue Gene?

Processor A's Memory

0xFFFFFFF

Processor B's Memory

0xFFFFFFF

Thread 3 stack

Heap

Thread 2 stack

Globals

Code

Heap

Globals

Code

0x00000000

0x00000000

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Aliasing Stack Data: Run Thread 2

Processor A's Memory

Processor B's Memory

0xFFFFFFF

Thread 2 stack

0xFFFFFFF

Thread 3 stack

Heap

Thread 2 stack

Globals

Code

Execution Copy

Heap

Globals

Code

0x00000000

0x00000000

53

Processor A's Memory

0xFFFFFFF

Processor B's Memory

0xFFFFFFF

Thread 3 stack

Heap

Thread 2 stack

Globals

Code

Heap

Globals

Code

0x0000000

0x00000000

Aliasing Stack Data: Run Thread 3

Processor A's Memory

Processor B's Memory

0xFFFFFFF

Thread 3 stack

0xFFFFFFF

Execution Copy

Thread 3 stack

Heap

Thread 2 stack

Globals

Code

Heap

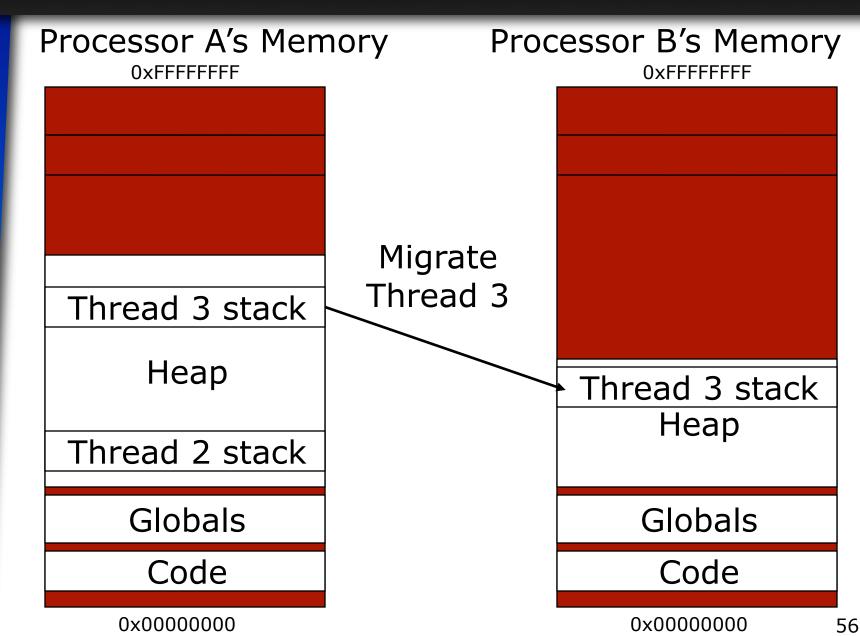
Globals

Code

0x00000000

0x00000000

55



Processor A's Memory

0xFFFFFFF

Processor B's Memory

0xFFFFFFF

Heap

Thread 2 stack

Globals

Code

Thread 3 stack Heap

Globals

Code

0x0000000

0x00000000

Processor A's Memory

Processor B's Memory

0xFFFFFFF

0xFFFFFFF

Execution Copy

Thread 3 stack

Heap

Thread 2 stack

Globals

Code

Thread 3 stack Heap

Globals

Code

0x00000000

0x00000000

- Does not depend on having large quantities of virtual address space
 - Works well on 32-bit machines
- Requires only one mmap'd region at a time
 - Works even on Blue Gene!
- Downsides:
 - Thread context switch requires munmap/ mmap (3us)
 - Can only have one thread running at a time (so no SMP's!)
- "-thread memoryalias" link time option

Heap Data

- Heap data is any dynamically allocated data
 - C "malloc" and "free"
 - C++ "new" and "delete"
 - F90 "ALLOCATE" and "DEALLOCATE"
- Arrays and linked data structures are almost always heap data

Migrate Heap Data

- Automatic solution: isomalloc all heap data just like stacks!
 - "-memory isomalloc" link option
 - Overrides malloc/free
 - No new application code needed
 - Same limitations as isomalloc; page allocation granularity (huge!)
- Manual solution: application moves its heap data
 - Need to be able to <u>size</u> message buffer, <u>pack</u> data into message, and <u>unpack</u> on other side
 - "pup" abstraction does all three

Delegation



Delegation

- Customized implementation of messaging
 - Enables Charm++ proxy messages to be forwarded to a delegation manager group
- Delegation manager
 - trap calls to proxy sends and apply optimizations
- Delegation manager must inherit from CkDelegateMgr class
- User program must to call
 - proxy.ckDelegate(mgrID);

Delegation Interface

```
.ci file
group MyDelegateMgr {
         entry MyDelegateMgr(); //Constructor
 };
.h file
class MyDelegateMgr : public CkDelegateMgr {
  MyDelegateMgr();
  void ArraySend(...,int ep,void *m,const
  CkArrayIndexMax &idx,CkArrayID a);
  void ArrayBroadcast(..);
  void ArraySectionSend(.., CkSectionID &s);
```

Array Multicast



Array Multicast/reduction library

- Array section a subset of chare array
- Array section creation
 - Enumerate array indices

```
    CkVec<CkArrayIndex3D> elems; // add array indices for (int i=0; i<10; i++) for (int j=0; j<20; j+=2) for (int k=0; k<30; k+=2) elems.push_back(CkArrayIndex3D(i, j, k)); CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, elems.getVec(), elems.size());</li>
```

- Alternatively, one can do the same thing by providing (lbound:ubound:stride) for each dimension:
 - CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, 0, 9, 1, 0, 19, 2, 0, 29, 2);
 - The above code creates a section proxy that contains array elements of [0:9, 0:19:2, 0:29:2].
- For user-defined array index other than CkArrayIndex1D to CkArrayIndex6D, one needs to use the generic array index type: CkArrayIndexMax.
 - CkArrayIndexMax *elems; // add array indices int numElems;
 CProxySection_Hello proxy = CProxySection_Hello::ckNew(helloArrayID, elems, numElems);

Array Section Multicast

- Once have the array section proxy
 - do multicast to all the section members:
 - CProxySection_Hello proxy;proxy.foo(msg) // multicast
 - send messages to one member using its local index
 - proxy[0].foo(msg)

Array Section Multicast

- Multicast via delegation
 - CkMulticast communication library
 - CProxySection_Hello sectProxy = CProxySection_Hello::ckNew(); CkGroupID mCastGrpId = CProxy_CkMulticastMgr::ckNew(); CkMulticastMgr *mcastGrp = CProxy_CkMulticastMgr (mCastGrpId).ckLocalBranch(); sectProxy.ckSectionDelegate(mCastGrpId); // initialize proxy sectProxy.foo(...); //multicast via delegation

Note, to use CkMulticast library, all multicast messages must inherit from CkMcastBaseMsg, as following:

```
class HiMsg : public CkMcastBaseMsg, public CMessage_HiMsg
{
    public:
    int *data;
};
```

Array Section Reduction

- Section reduction with delegation
- Use default reduction callback

```
CProxySection_Hello sectProxy;
```

CkMulticastMgr *mcastGrp = CProxy_CkMulticastMgr (mCastGrpId).ckLocalBranch();

mcastGrp->setReductionClient(sectProxy, new CkCallback(...));

Reduction

```
CkGetSectionInfo(sid, msg);
CkGetIback cb(CkIndex, myArray::foo(NIIII I ) this
```

CkCallback cb(CkIndex_myArray::foo(NULL),thisProxy); mcastGrp->contribute(sizeof(int), &data, CkReduction::sum_int, sid, cb);

With Migration

- Works with migration
 - When intermediate nodes migrate
 - When migrates, multicast tree will be automatically rebuilt
 - Root processor
 - Application needs to initiate the rebuild
 - Will change to automatic in future

SDAG



Structured Dagger

- What is it?
 - A coordination language built on top of Charm++
 - Express control flow in interface file
- Motivation
 - Charm++'s asynchrony is efficient and reliable, but tough to program
 - Split phase Flags, buffering, out-of-order receives, etc.
 - Threads are easy to program, but less efficient and less reliable
 - Implementation complexity
 - Porting headaches
 - Want benefits of both!

Structured Dagger Constructs

- when <method list> {code}
 - Do not continue until method is called
 - Internally generates flags, checks, etc.
 - Does not use threads
- atomic {code}
 - Call ordinary sequential C++ code
- if/else/for/while
 - C-like control flow
- overlap {code1 code2 ...}
 - Execute code segments in parallel
- forall
 - "Parallel Do"
 - Like a parameterized overlap

Stencil Example Using SDAG

```
array[1D] myArray {
entry void GetMessages () {
    when rightmsgEntry(), leftmsgEntry() {
         atomic {
           CkPrintf("Got both left and right messages \n");
           doWork(right, left); }
     }
};
entry void rightmsgEntry();
entry void leftmsgEntry();
```

Overlap for LeanMD Initialization

```
array[1D] myArray {
 entry void waitForInit(void) {
     overlap {
        when recvNumCellPairs(myMsg* pMsg) {
          atomic { setNumCellPairs(pMsg->intVal); delete
pMsg; }
        when recvNumCells(myMsg * cMsg) {
          atomic { setNumCells(cMsg->intVal); delete cMsg; }
     }
```

For for LeanMD timeloop

```
entry void doTimeloop(void) {
   for (timeStep_=1; timeStep_<=SimParam.NumSteps; timeStep++) {</pre>
        atomic {sendAtomPos(); }
        overlap {
          for (forceCount_=0; forceCount_<numForceMsg_; forceCount_++) {</pre>
            when recvForces(ForcesMsg* msg) { atomic {procForces(msg); } }
          for (pmeCount_=0; pmeCount_<nPME; pmeCount_++) {</pre>
            when recvPME(PMEGridMsg* m) {atomic {procPME(m);}}
        atomic { doIntegration(); }
        if (timeForMigrate()) { ... }
```

Thank You!

Free source, binaries, manuals, and more information at:

http://charm.cs.uiuc.edu/

Parallel Programming Lab at University of Illinois

